

## Communications IV Exercise Solutions

### Exercise 6.1

$$\begin{aligned}
 H(x) &= -\sum_{i=1}^6 P(i) \ln P(i) \text{ nits} \\
 &= -\{0.1\ln(0.1) + 0.2\ln(0.2) + 0.3\ln(0.3) + 0.05\ln(0.05) + 0.15\ln(0.15) + 0.1\ln(0.1)\} \\
 &= 1.670 \text{ nits} = 2.409 \text{ bits} \\
 H(\text{uniform}) &= \log_2(6) = 2.585 \text{ bits}
 \end{aligned}$$

### Exercise 6.7

The line  $x-1$  is tangent to  $\ln(x)$  at  $x = 1$  and a simple sketch shows that  $\ln(x) \leq (x-1)$  for all  $x > 0$ .

$$\begin{aligned}
 H(x) &= -\sum_{i=1}^N P_i \ln P_i \\
 H(y) &= \ln N = \sum_{i=1}^N P_i \ln N \\
 H(x) - H(y) &= \sum_{i=1}^N P_i \ln (1/NP_i) \\
 &\leq \sum_{i=1}^N P_i [1/NP_i - 1] = 1 - 1 = 0 \\
 H(x) &\leq H(y)
 \end{aligned}$$

### Exercise 6.22

1.

$$\begin{aligned}
 H(x) &= -\sum_{i=1}^4 P(i) \ln P(i) \text{ nits} \\
 &= -\{0.1\ln(0.1) + 0.2\ln(0.2) + 0.3\ln(0.3) + 0.4\ln(0.4)\} \\
 &= 1.280 \text{ nits} = 1.847 \text{ bits}
 \end{aligned}$$

2. Average code length  $\geq 1.847$  bits.

3.

		Code	$n_i P_i$
A	0.4	0	0.40
B	0.3	0 1.0	10    0.60
C	0.2	0    0.3 1	110    0.60
D	0.1	1	111 <u>0.30</u> <u>1.90</u>

Average code word length = 1.90 bits (Efficiency = 97.2%).

## Exercise 6.56

1. This is trivial.
2. If  $P_b$  is the probability of a bit error,

$$MSE_{chan} = P_b \Delta^2 \left(1 + 4 + 4^2 + \dots + 4^{v-1}\right) = P_b \Delta^2 \frac{4^v - 1}{3}$$

3. With  $MSE_{quant} = \Delta^2/12$  and defining  $x_{max} = N\Delta/2$  gives the expression required.

$$MSE = \frac{x_{max}^2}{3N^2} \left[1 + 4P_b(4^v - 1)\right]$$

You can also deduce that the channel MSE and quantisation MSE are equal when

$$P_b = \frac{1}{4(4^v - 1)} = 3.81 \times 10^{-6} \text{ for } v=8.$$

This corresponds to a SNR of 13.0 dB with gaussian noise.

4. Simply calculate the ratio of the signal and noise powers.

$$\text{Signal} = x(t), \quad \bar{x}(t) = \frac{x(t)}{x_{max}} \text{ [this would be } m(t) \text{ in our notation].}$$

$$\text{SNR} = \frac{3N^2 \langle x^2 \rangle}{x_{max}^2 [1 + 4P_b(4^v - 1)]} = \frac{3N^2 \langle \bar{x}^2 \rangle}{1 + 4P_b(4^v - 1)}$$

## Solutions to Exercises on Slides

### Slide 6.07

For a source  $x$  with  $n$  symbols the probabilities are  $P_i$ ,  $i = 1, 2, \dots, n$ . Since the sum of these must be unity, we have only  $n-1$  independent variables.

$$\begin{aligned}\sum_{i=1}^n P_i &= 1 \Rightarrow P_n = 1 - \sum_{i=1}^{n-1} P_i \\ H(x) &= -\sum_{i=1}^n P_i \ln(P_i) \text{ nits} \\ \frac{\partial H(x)}{\partial P_i} &= -[1 + \ln(P_i)] - [1 + \ln(P_n)] \frac{\partial P_n}{\partial P_i} ; i = 1, 2, \dots, n-1 \\ &= -[1 + \ln(P_i)] + [1 + \ln(P_n)] \quad \text{since } \frac{\partial P_n}{\partial P_i} = -1 \\ &= \ln\left(\frac{P_n}{P_i}\right) = 0 \Rightarrow P_i = P_n = \frac{1}{n} ; i = 1, 2, \dots, n-1 \\ H_{\max}(x) &= -\sum_{i=1}^n \frac{1}{n} \ln\left(\frac{1}{n}\right) = \ln(n) \text{ nits} = \log_2(n) \text{ bits}\end{aligned}$$